

Computation and Control

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Computation and reaction

- A reactive system is a system that describes how to react to events
 - “moving left” → “steer right”
 - “slowing down” → “increase thrust”
- Algorithms are a set of rules to apply repeatedly
 - $while(i \neq 0) \{ k \ast = i; i-- \}$



Rewriting systems

- A *rewriting* system uses a language \mathcal{L} , together with a set of rules $s_i \rightarrow t_i$ for some $s_i, t_i \in \mathcal{L}$.
- s_i is called a *redex* (a pattern or set of terms)
- t_i is called a *contractum*
- A *computation* is a sequence of rewrite applications $e_1 \rightarrow e_2 \rightarrow \dots \rightarrow e_n$



Lambda calculus (canonical example)

Language \mathcal{L} :

e	::=	v	variables
		$e_1 e_2$	function application
		$\lambda v.e$	function abstraction

Rewrites:

$$(\lambda v.e_1) e_2 \rightarrow_{\beta} e_1[e_2/v]$$

	$(\lambda x.\lambda y.x + y) 1 2$
\rightarrow	$(\lambda y.1 + y) 2$
\rightarrow	$1 + 2$
\rightarrow	3

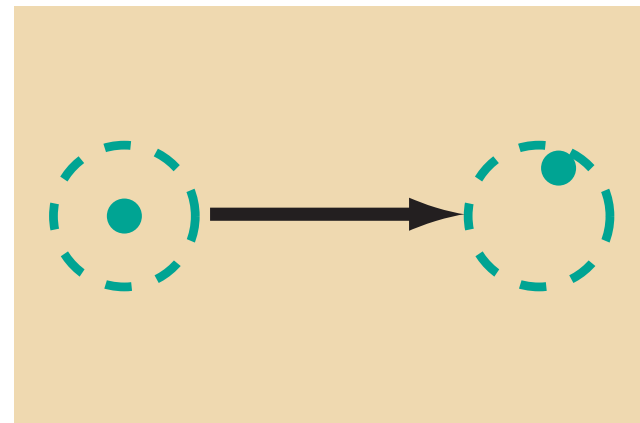
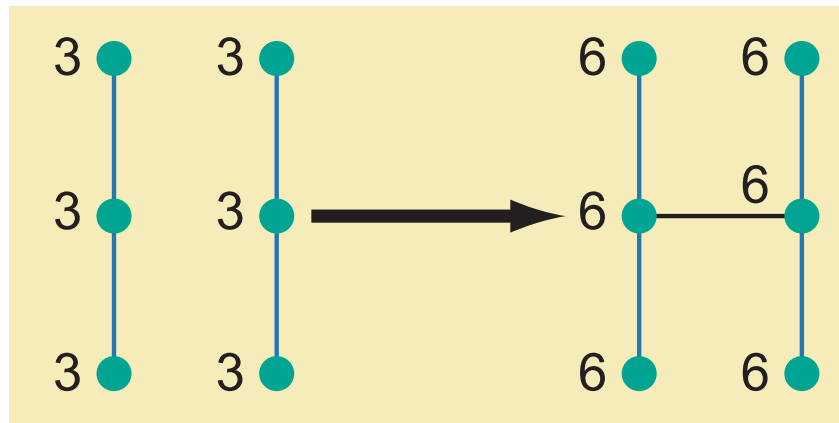
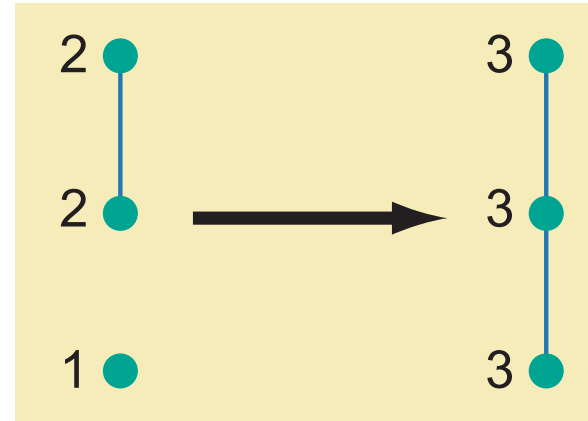
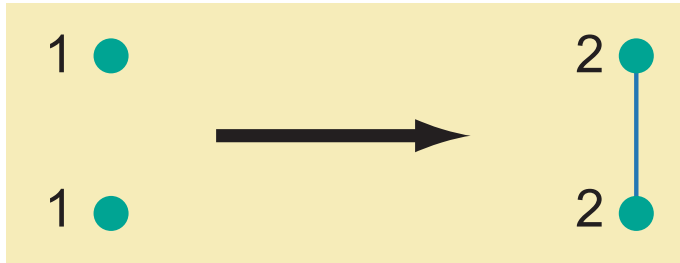


Example: vehicle formation/assembly

- Assemble a set of vehicle/parts into a formation



Formation rules



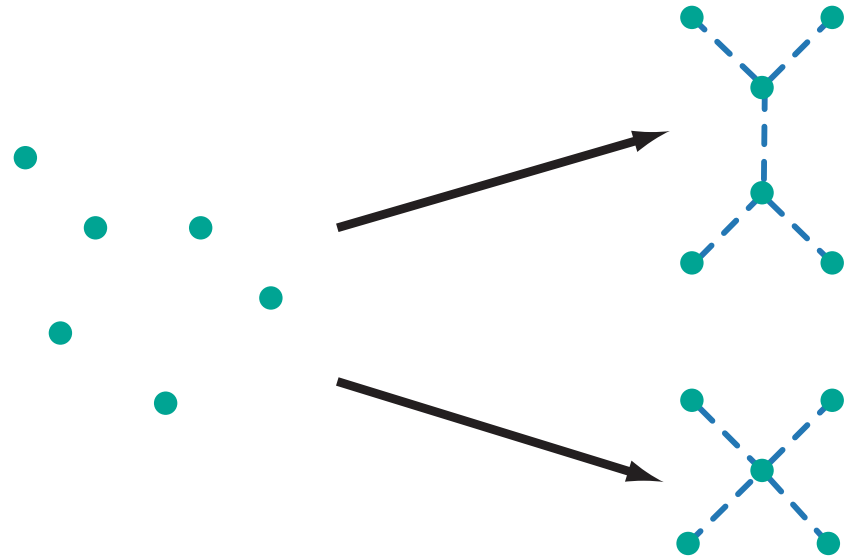
Rewriting rules

- Rule classes
 - Progress: getting closer to a goal
 - Dynamics
 - Adversaries: destructive actions
- Technical questions
 - *Determinism (Church-Rosser)*
 - *Progress (liveness)*
 - *Termination*
 - *Locality*



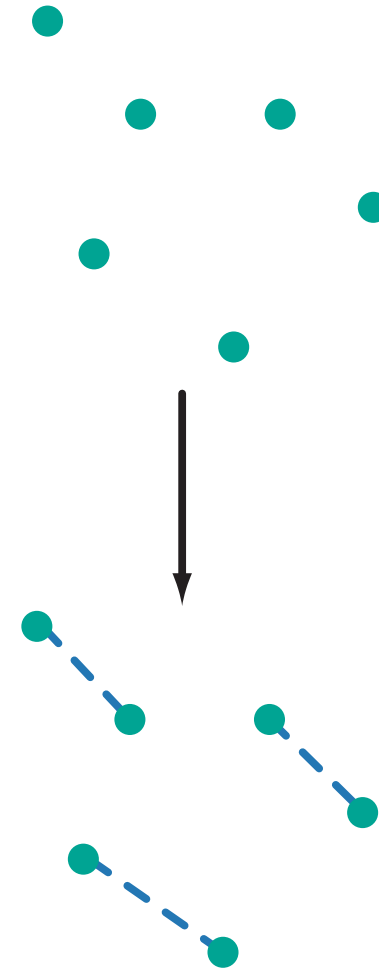
Determinism (Church-Rosser)

- Does the order of evaluation matter?
- Often the answer is no
 - *There may be reasons to have more than one result*
- Proofs are quite difficult



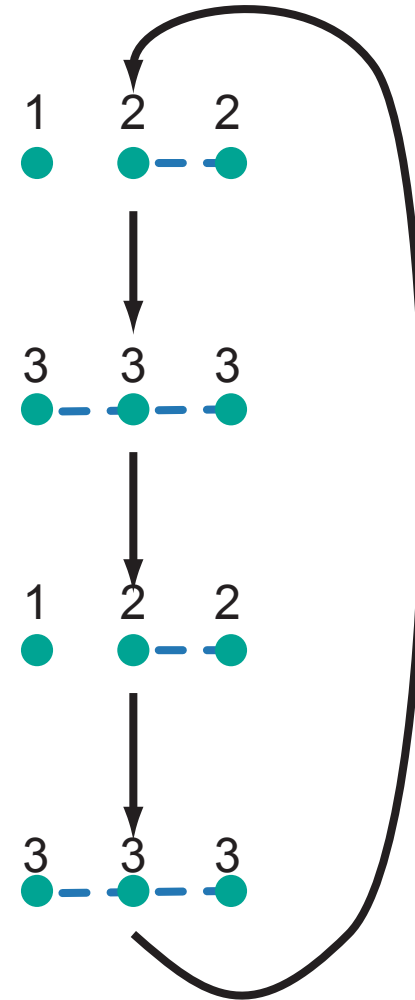
Deadlock/progress

- Is it possible to build a partial, final, formation?
 - *Show that: if a formation is not final, at least one rule is always enabled*
- Easy to prevent
 - *Add more rules to make progress from partial formations*
 - *Add “undo” to reverse bogus computations*



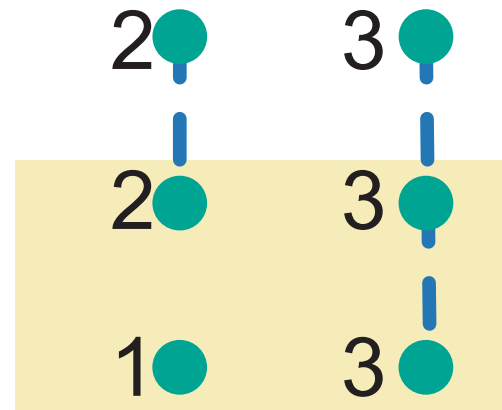
Termination

- Does every reduction sequence terminate?
- Termination proofs require the formation of a metric
 - *Often infeasible, if not it can be extremely hard*
 - *In many cases it doesn't matter*



Locality

- Is it possible to make decisions locally?
 - *Locality is determined by the scope of the rewrite*
- Methods
 - *Optimize the program to limit the scope*
 - *Introduce communication*



Languages

- Rewriting languages

- *Determinism*(*)
- *Termination*(*)
- *Progress*
- *Locality*

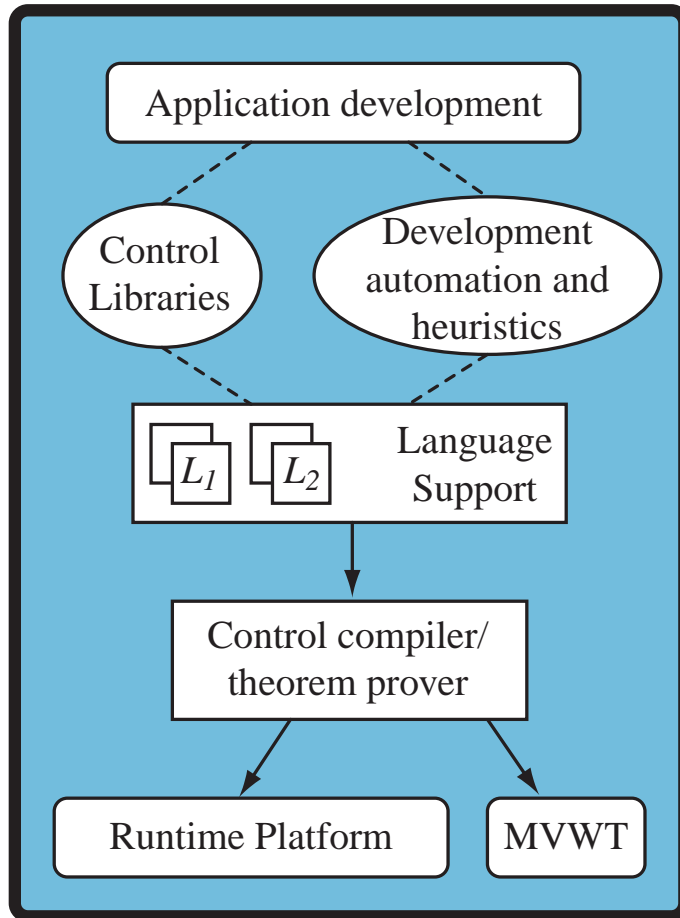
- (*) *hard to prove, not always useful*

UNITY-like languages:

G_1	\rightarrow	P_1	decisions and actions
G_2	\rightarrow	P_2	
		\vdots	
G_n	\rightarrow	P_n	
<hr/>			
G_c	\rightarrow	P_c	physics, control
G_a	\rightarrow	P_a	adversaries



Logical Programming Environments

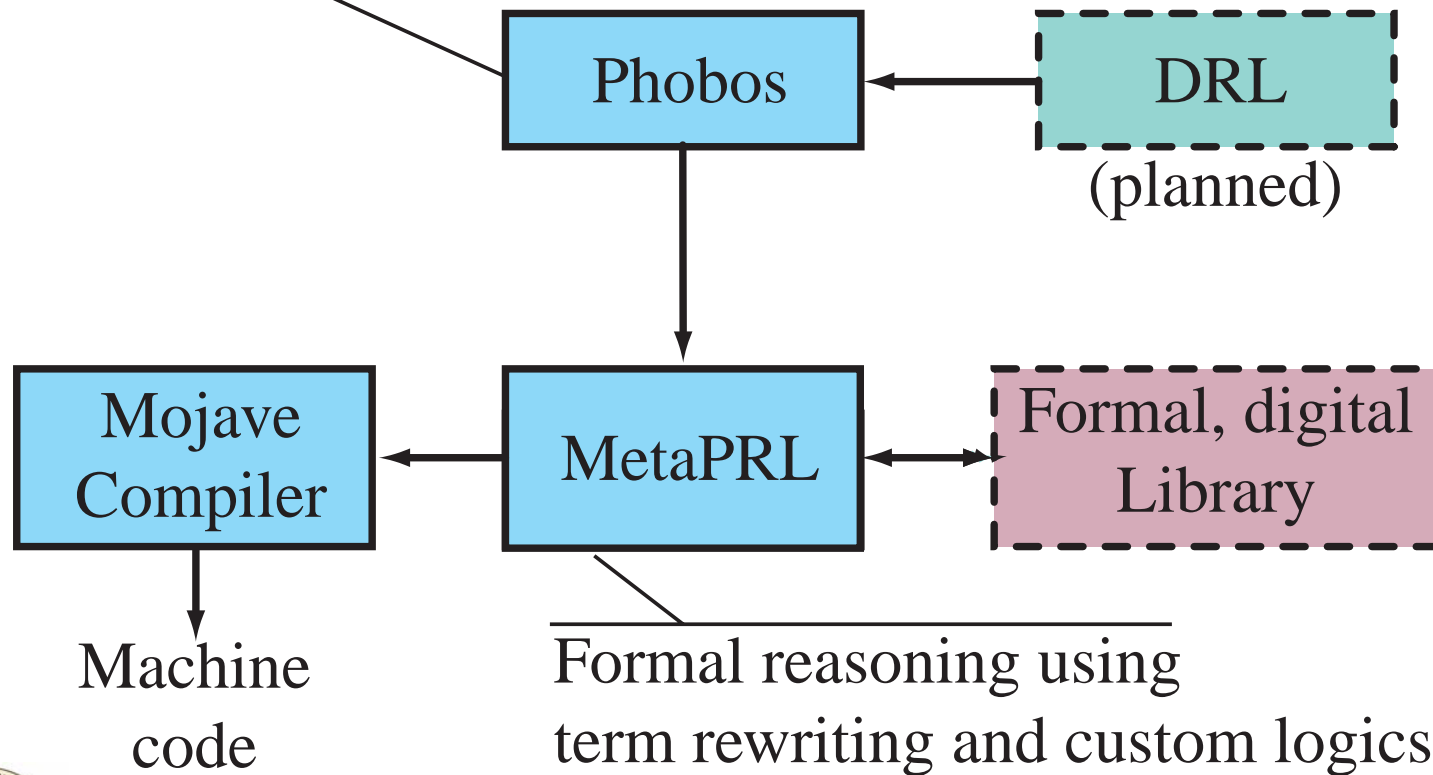


- The LPE is a framework for supporting formal design
 - *Type theory is a common language for specification and synthesis*
 - *Enables collaborative development of verified control libraries and design automation tools*
 - *The compiler is an assistant, and the link to executable code*



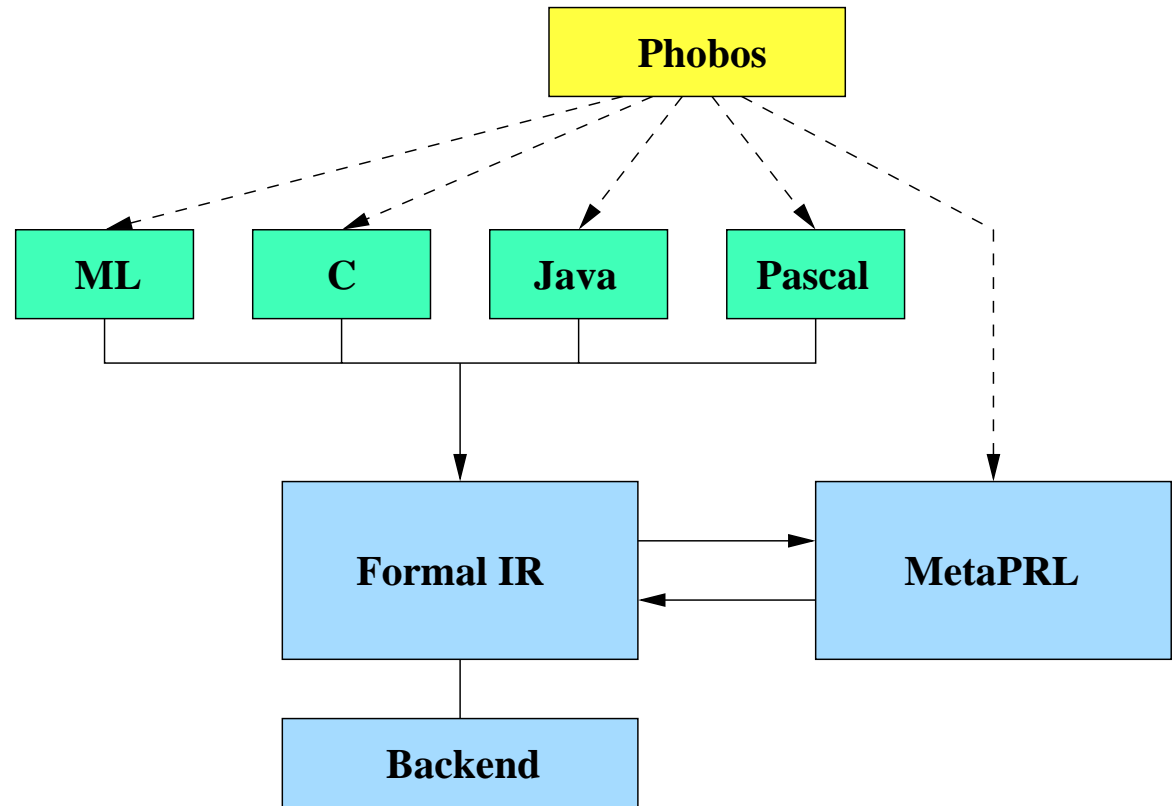
Logical Programming Environments

Definitions of languages,
syntax, rewrite rules



Phobos

- Phobos is a front-end for domain-specific languages
 - *Programs are translated to one of a set of “standard” languages*
 - *or to a theorem prover*



Language definitions

- Each language has a lexicon

```
Tokens -longest {  
  NUM = "[0-9]+"      { __token__[p:s]{'pos} -> num[p:s]{'pos} }  
  ...  
  * SPACE = " "      {}  
}
```

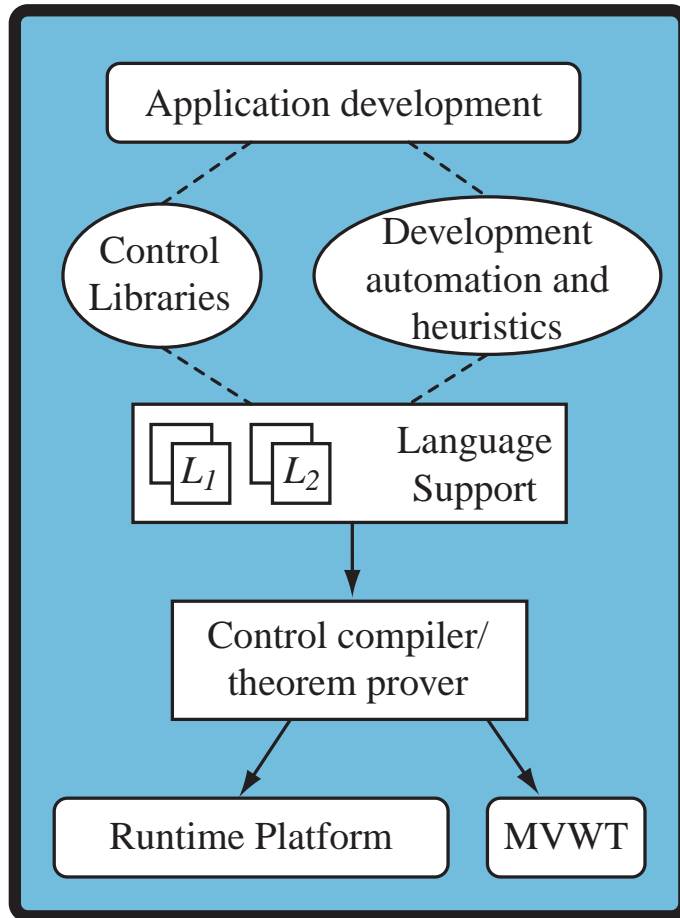
- And a grammar

```
%left PLUS MINUS  
%left TIMES DIV  
%left LPAREN RPAREN
```

```
Grammar -start exp {  
  exp ::= NUM      { num[p:s]{'pos} -> exp{num[p:s]; 'pos}}  
  | ID             { ... }  
  | exp PLUS exp  { 'e1 PLUS 'e2 -> exp{sum{'e1; 'e2};  
                                     union_exp_pos{'e1; 'e2} }  
  ...  
}
```



Logical Programming Environments



• Parts

- *Phobos language support complete*
- *DRL language design and control primitives*
- *Code compiler complete*

